

NETWORK QUALITY OF SERVICE ON BIG SCALE

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NATIONAL RESEARCH, DEVELOPMENT
AND INNOVATION OFFICE
HUNGARY

PROGRAM
FINANCED FROM
THE NRDI FUND

Results (2022 aug. – 2023 jan.)

- **Accepted Paper**

- **Gergő Gombos**, Dávid Kis, Lilla Tóthmérész, Tamás Király, Szilveszter Nádas, and Sándor Laki: „**Flow Fairness with Core-Stateless Resource Sharing in Arbitrary Topology**”, IEEE/ACM Transactions on Networking, 2022 (Q1)

- **Presented Demo**

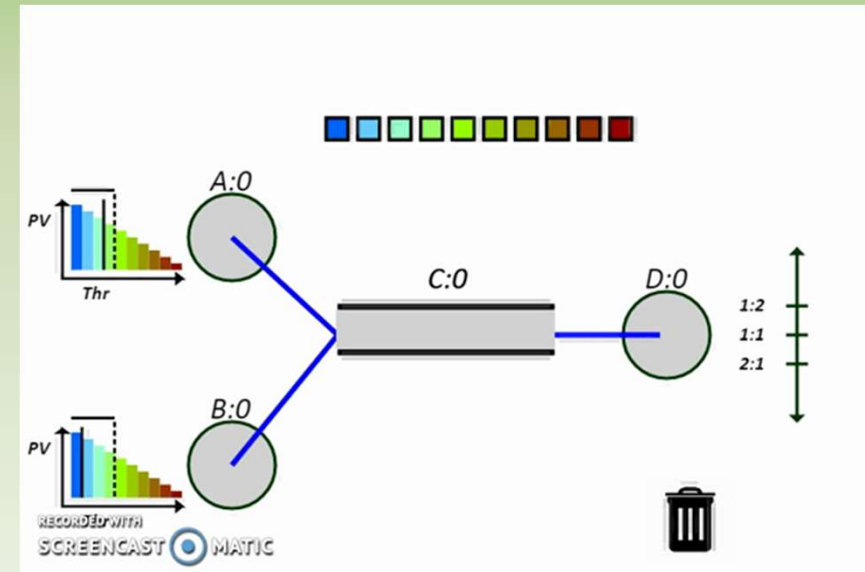
- Dávid Kis, **Gergő Gombos**, Szilveszter Nádas, and Sándor Laki: „**Resource Sharing Beyond FQ: 35K Users at 100Gbps**”, IEEE/ACM SIGCOMM Demo, 2022 (A*)

- **Ongoing Project**

- Network Hierarchical Quality of Service

Reminder: Core-Stateless Resource Sharing Control with PPV

- PPV is a Core-Stateless Resource Sharing solution
- Marking at the network edge
 - every flow has an own policy
 - the packet marked based on this policy
- AQM
 - decision based on the packet values only
 - no information about the
 - policies
 - flows
 - fast and easy implementation



Tutorial video @ ppv.elte.hu

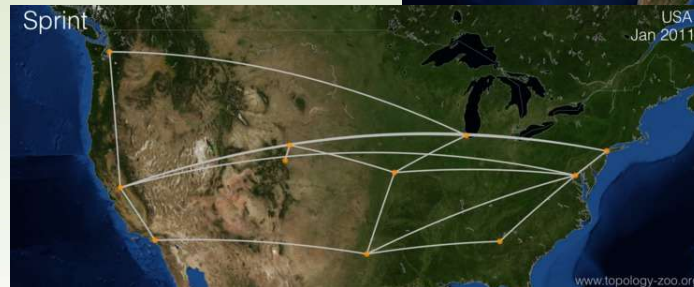
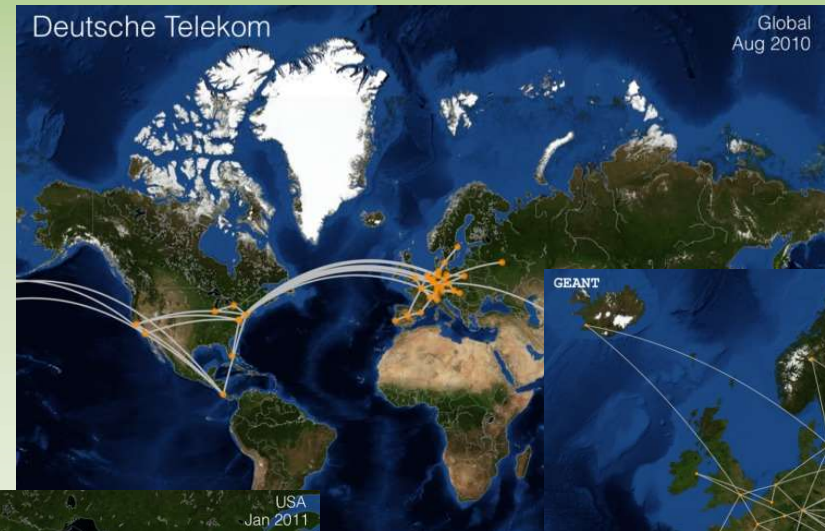
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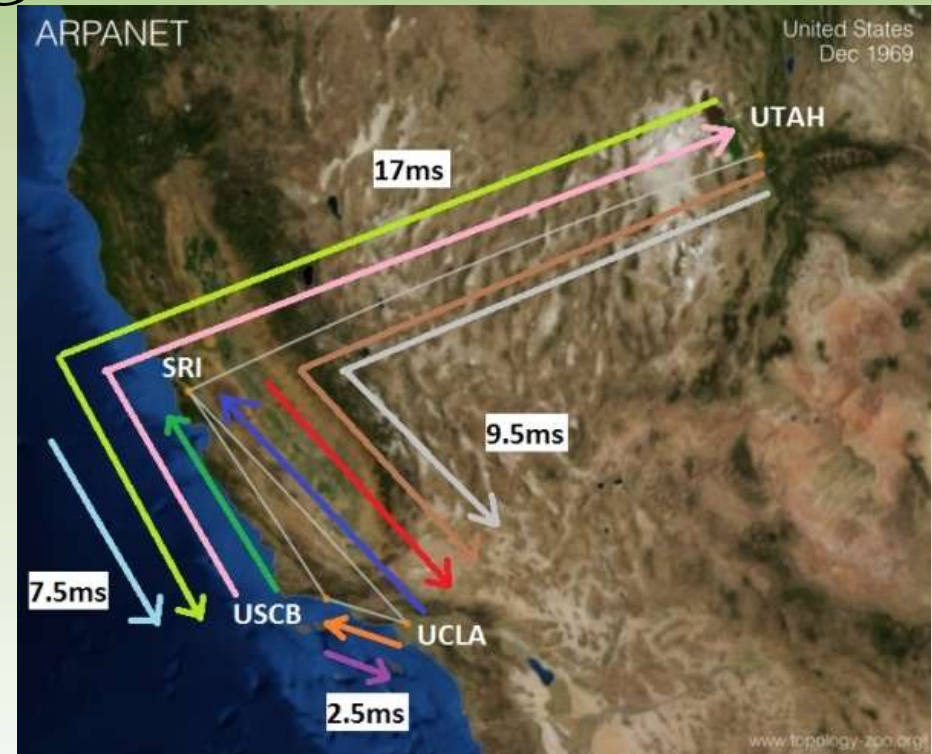
PPV Core-Stateless Resource Sharing Control with PPV on large networks

- topology-zoo.org
 - real network topologies
 - nodes, edges, coordinates
 - it does not have RTTs
- Used topologies:
 - Sprint
 - Deutsche Telekom
 - Geant



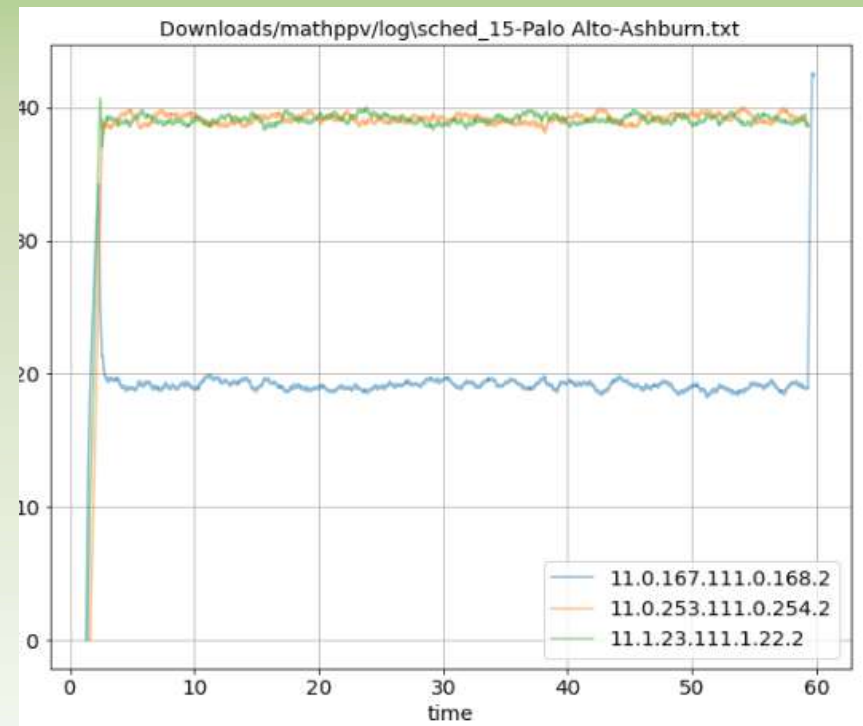
PPV Core-Stateless Resource Sharing Control with PPV on large networks

- Used protocols:
 - DCTCP
 - works well with small RTTs
 - NewReno
 - half the rate when dropping is
 - UDP
 - flow model simulation, no congestion control
- Simulation:
 - random $n^2/10$ flow (DT, Geant)
 - random n^2 flow (Sprint)
 - 50% Silver, 50% Gold



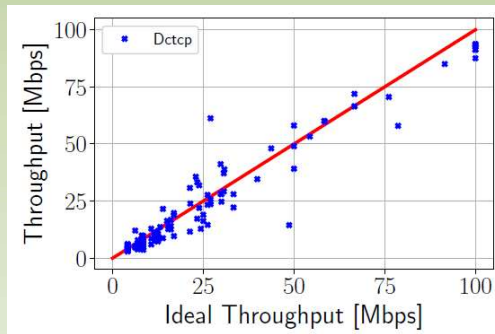
PPV Core-Stateless Resource Sharing Control with PPV on large networks

- 3 Silver flow on one node
- Blue lower, because it is reduce they rate earlier
- Others use the available bandwidth

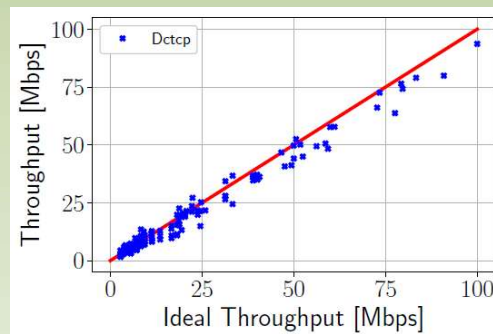


PPV Core-Stateless Resource Sharing Control with PPV on large networks

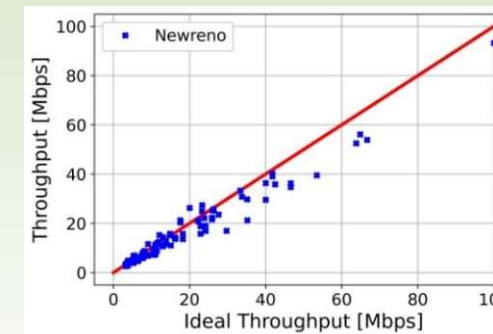
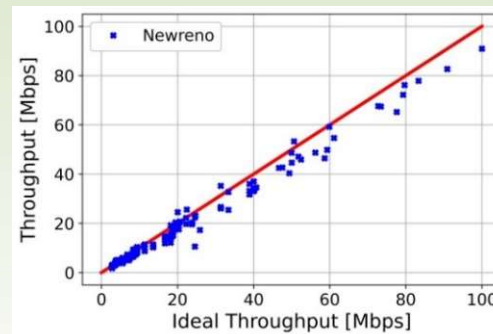
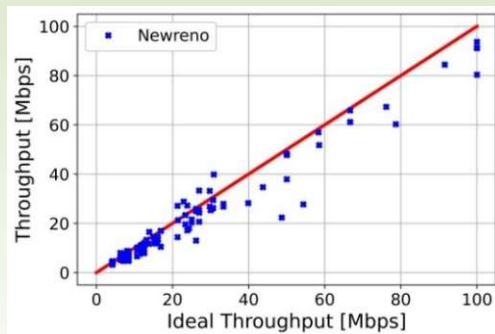
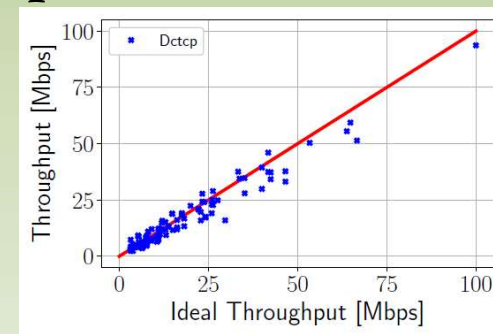
Deutsche Telekom



Geant



Sprint

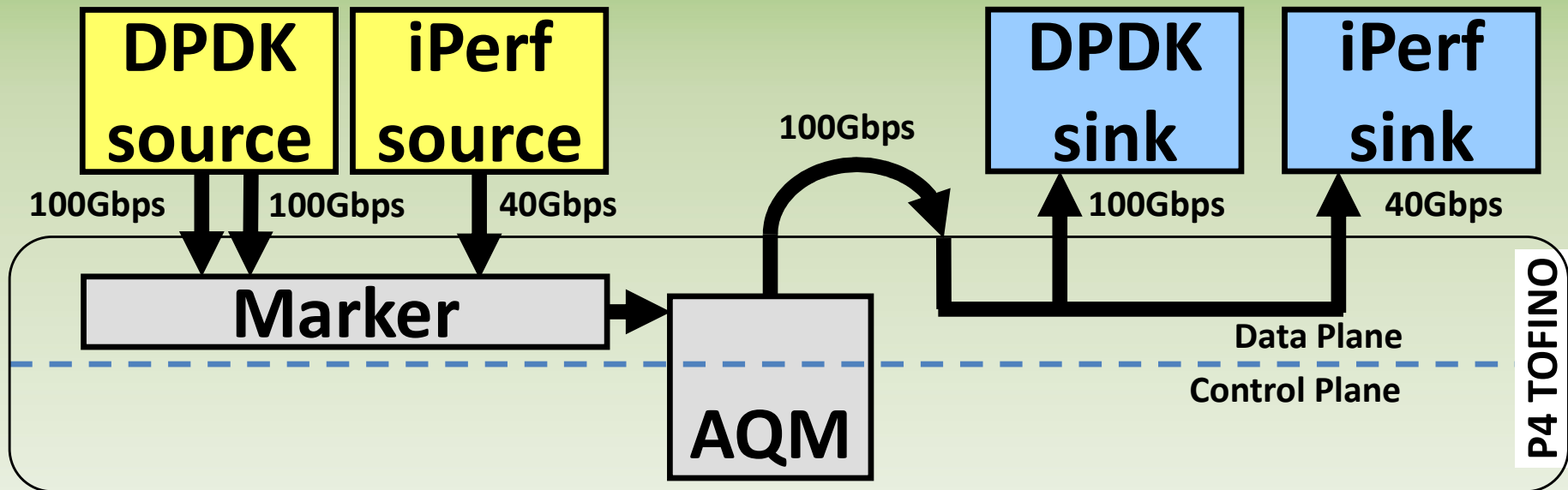


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- **Presented Demo**

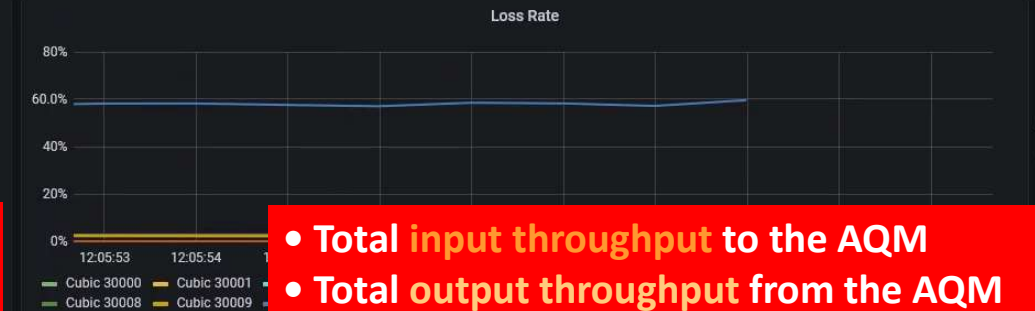
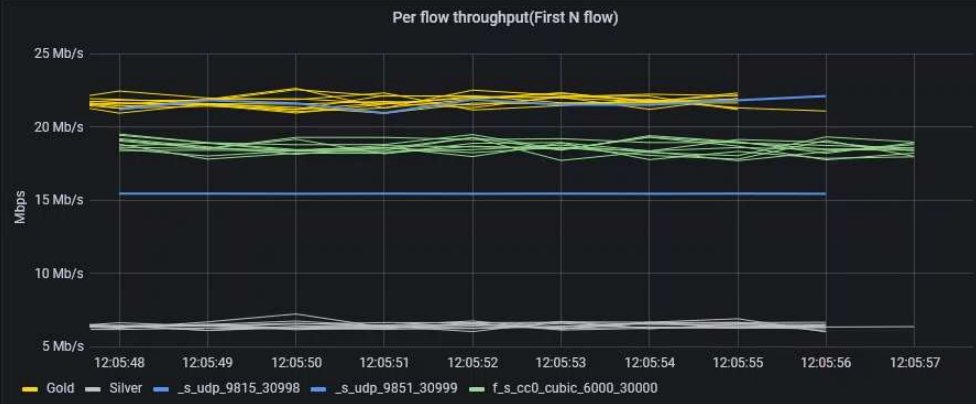
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Testbed



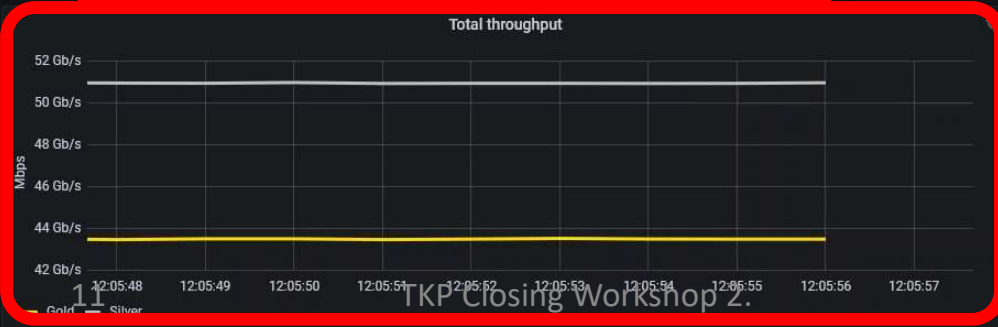
Grafana dashboard

Gold portrange min 20000 Gold portrange max 21999 Silver portrange min 22000 Silver portrange max 29999 Flow count 10



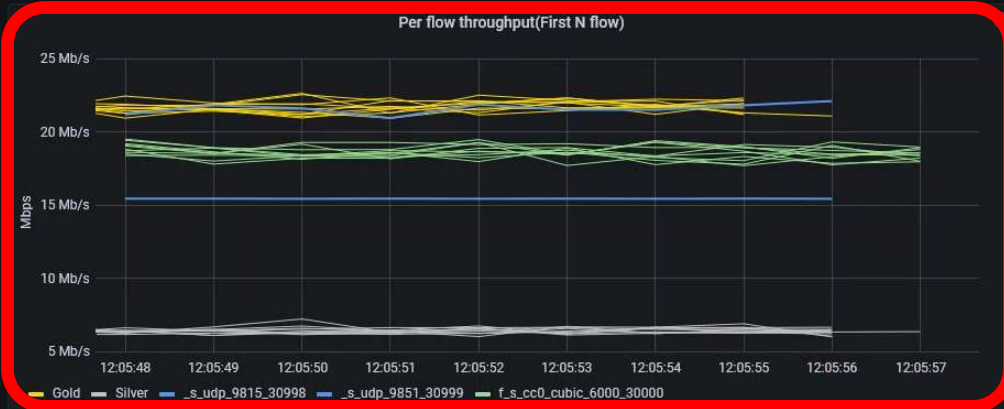
Gold flows
2000

Silver flows
8000



TKP Closing Workshop 2.

Gold portrange min 20000 Gold portrange max 21999 Silver portrange min 22000 Silver portrange max 29999 Flow count 10

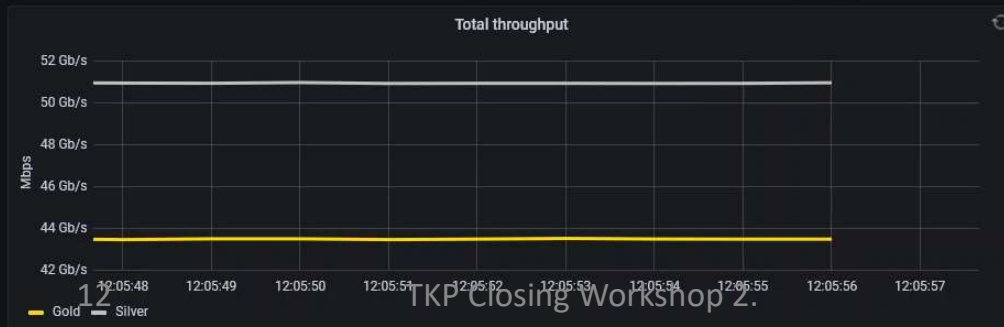


- 10 selected Gold UDP flows via DPKD
- 10 selected Silver UDP flows via DPKD
- 2 designated UDP flows via iPerf
- 10 Cubic TCP flows via iPerf

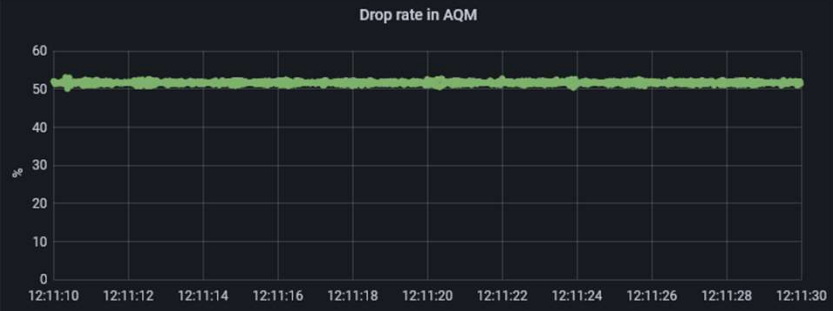
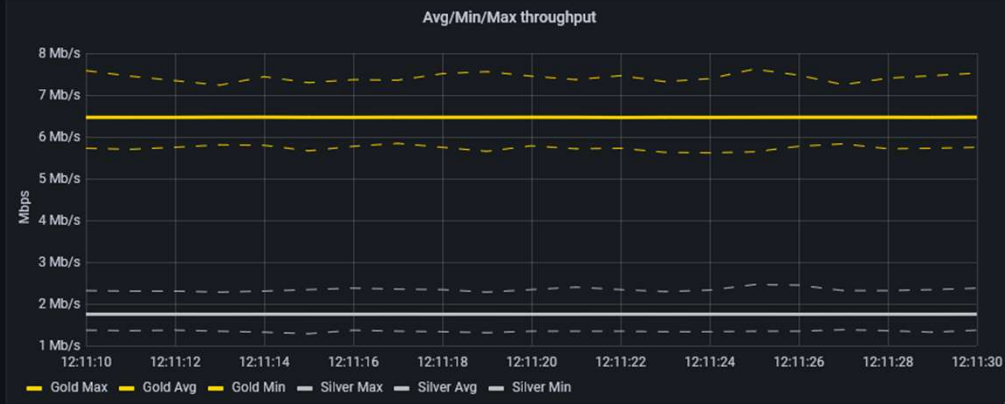
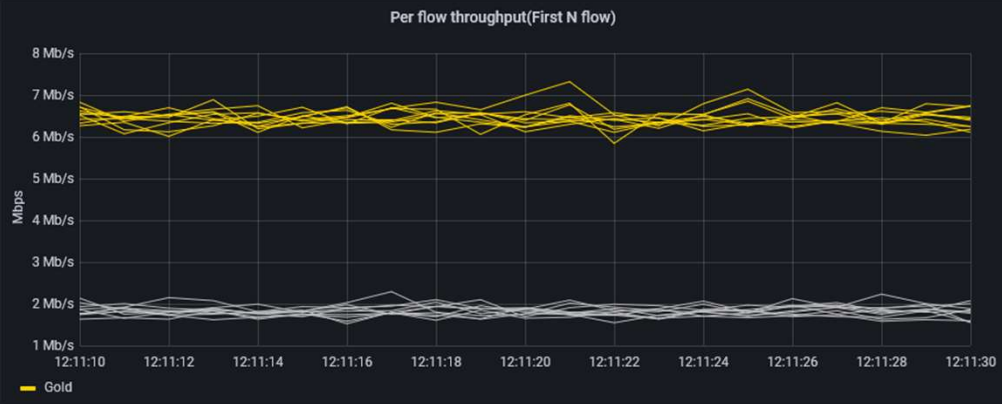
Average/min/max throughput of UDP flows generated by DPKD pktgen



Silver flows
8000

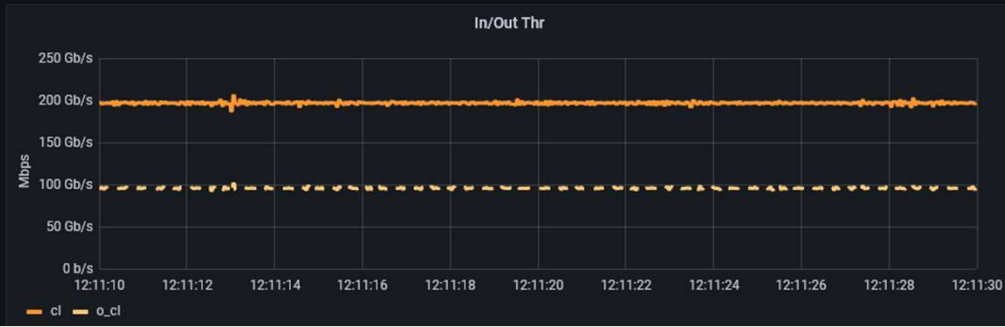
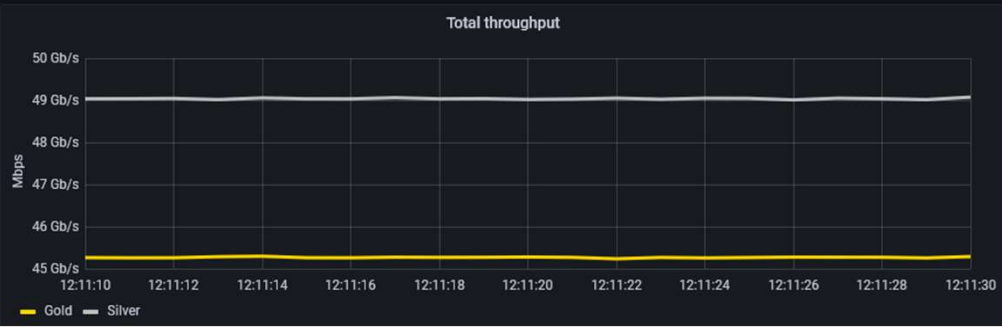


Gold portrange min 20000 Gold portrange max 26999 Silver portrange min 27000 Silver portrange max 60000 Flow count 10

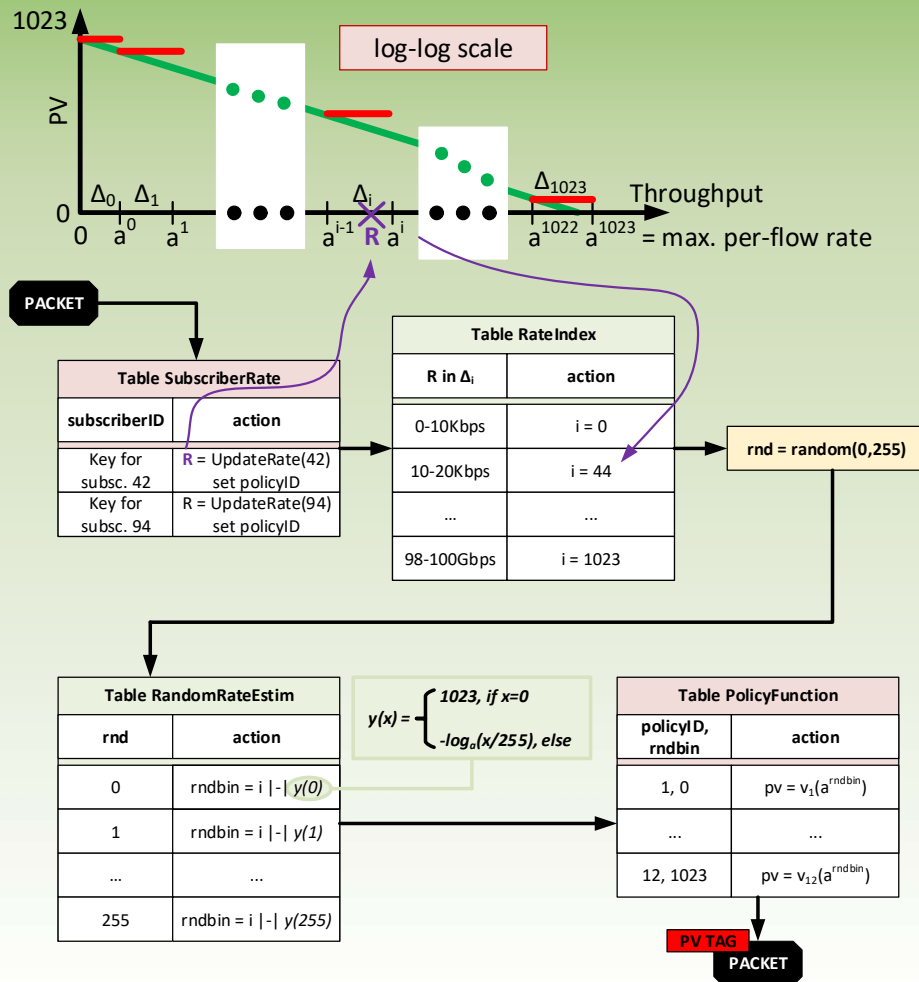


Gold flows
7000

Silver flows
28000



Packet Marking on P4 Tofino



- We used the built-in rate measurement and convert it to index
- With the index and a random number we calculate a random rate
- The policy function is quantized logarithmically and we can select the PV based on the random rate
- Data plane only implementation.
 - Policy functions are configured by the control plane
- We run 35000 marker instances in parallel

Results (2022 aug. – 2023 jan.)

- **Ongoing Project**
 - Network Hierarchical Quality of Service

Mobile backhaul

- Traffic aggregates instead of flows
- Resource sharing hierarchy
- Physical network shared among multiple operators

Infrastructure-side expectation

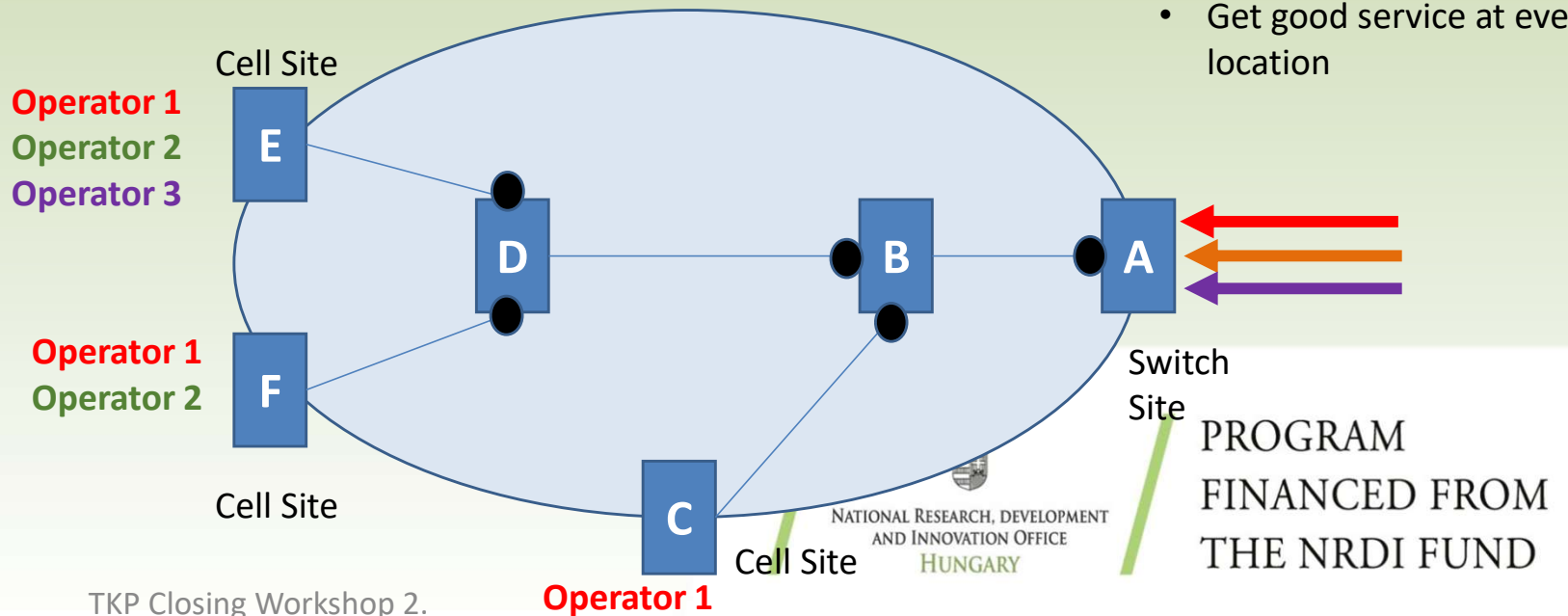
- Utilizing resources and serving operators' needs

Operator-side expectation

- Delivering traffic between the switch site and the cells where its subscribers are located
- Getting its deserved share according to its SLA/payment

User-side expectation

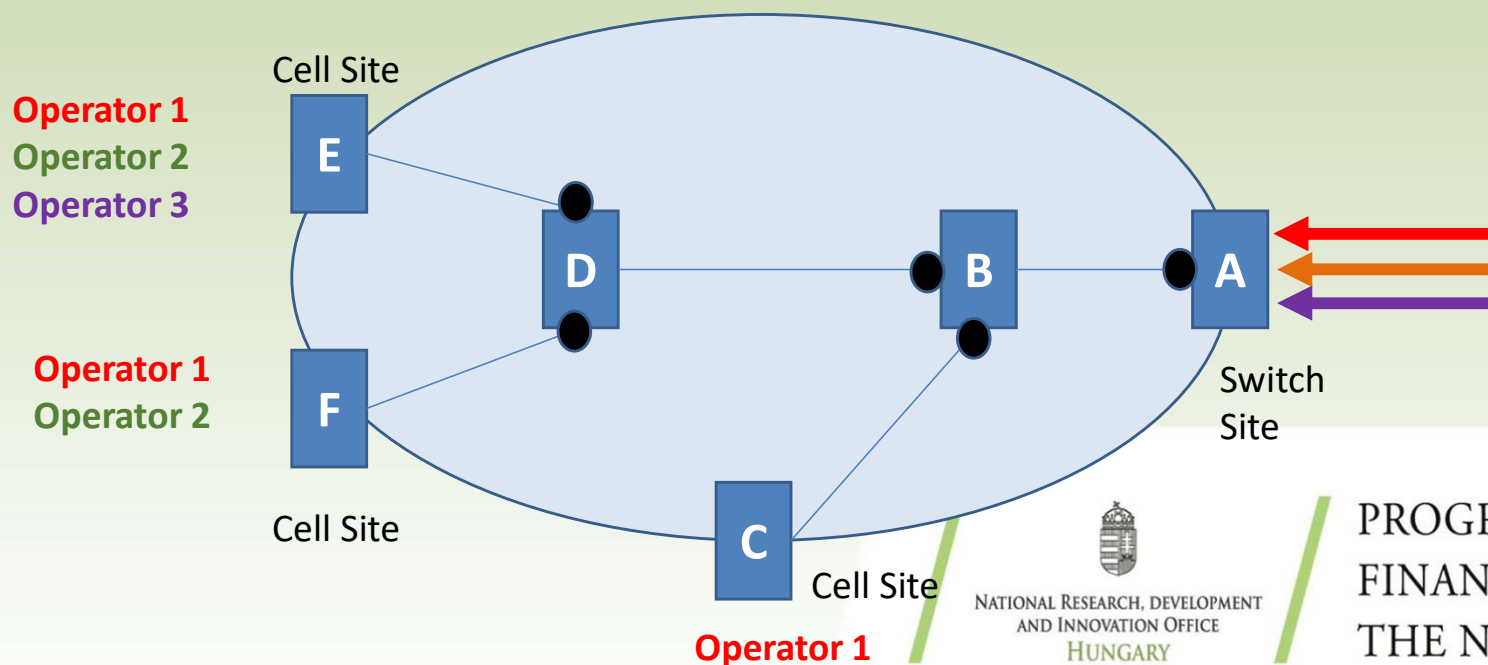
- Get good service at every location



Current approach Link-by-link HQoS

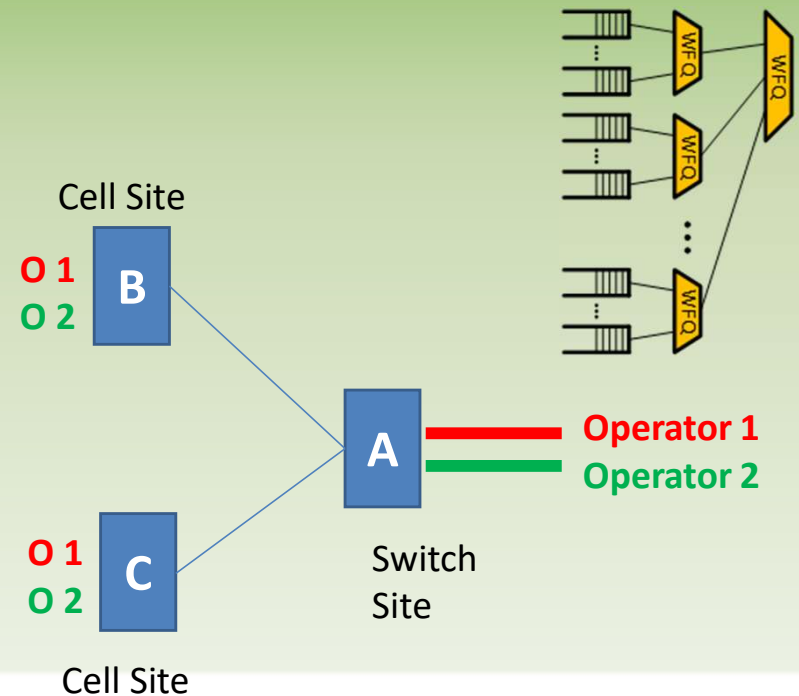
Operator 1 = VLAN 100 -> weight 40%
Operator 2 = VLAN 200 -> weight 30%
Operator 3 = VLAN 300 -> weight 30%

● Egress radiolink interfaces

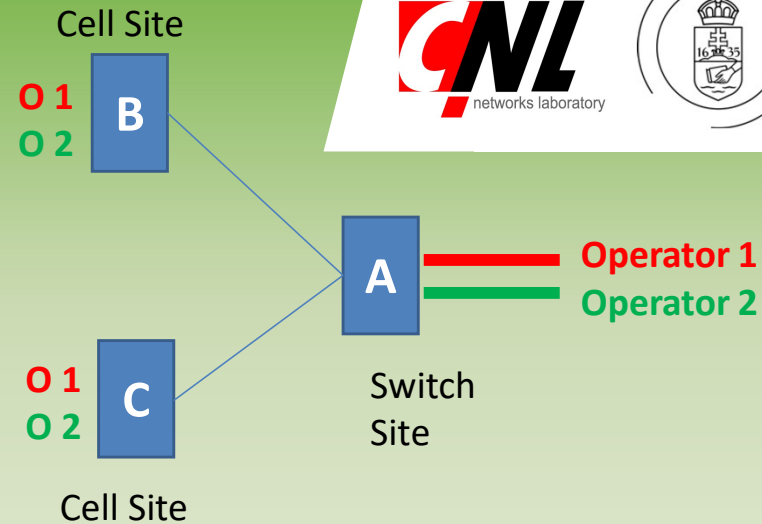


Simple topology – downstream case

- Two bottleneck links
 - A-B 100Mbps
 - A-C 100Mbps
- Two operators
 - Both can send towards cell B and C
- HQoS marking at site A
 - Operator 1 and 2 has 1:1 share
 - Users within each operator are equal.



Link-by-link VS Network-wide HQoS

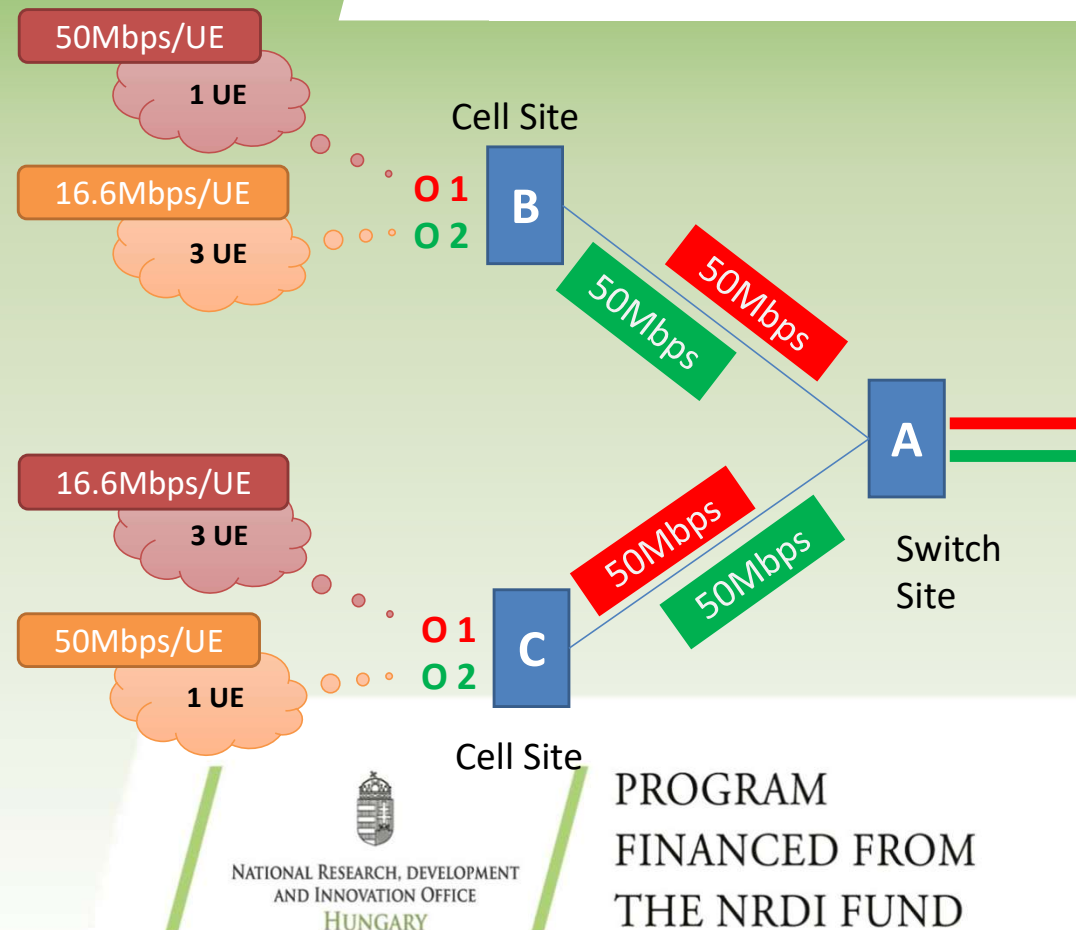


Case	Traffic				Network-Wide HQoS (NHQoS)										Link-by-link HQoS (LHQoS)									
	O1		O2		O1 throughput [Mbps]					O2 throughput [Mbps]					O1 throughput [Mbps]					O2 throughput [Mbps]				
	#users		#users		per user		total			per user		total			per user		total			per user		total		
	AB	AC	AB	AC	AB	AC	AB	AC	total	AB	AC	AB	AC	total	AB	AC	AB	AC	total	AB	AC	AB	AC	total
1	1		4		50		50		50	12.5		50		50	50		50		50	12.5		50		50
2	1	3	3	1	25	25	25	75	100	25	25	75	25	100	50	16.6	50	50	100	16.6	50	50	50	100
3	1	3	1	3	50	16.6	50	50	100	50	16.6	50	50	100	50	16.6	50	50	100	50	16.6	50	50	100
4	1	3	4		20	33.3	20	100	120	20		80		80	50	33.3	50	100	150	12.5		50		50
5	3	1	4		14.3	100	42.9	100	142.9	14.3		57.1		57.1	16.6	100	50	100	150	12.5		50		50
6	1	3	2	2	40	20	40	60	100	30	20	60	40	100	50	16.6	50	50	100	25	25	50	50	100
7	1	3	20	20	40	20	40	60	100	3	2	60	40	100	50	16.6	50	50	100	2.5	2.5	50	50	100

Table 1: Resource sharing examples. The minimum per-user throughput (marked with boldface) and the user fairness within each MNO are improved in most cases.

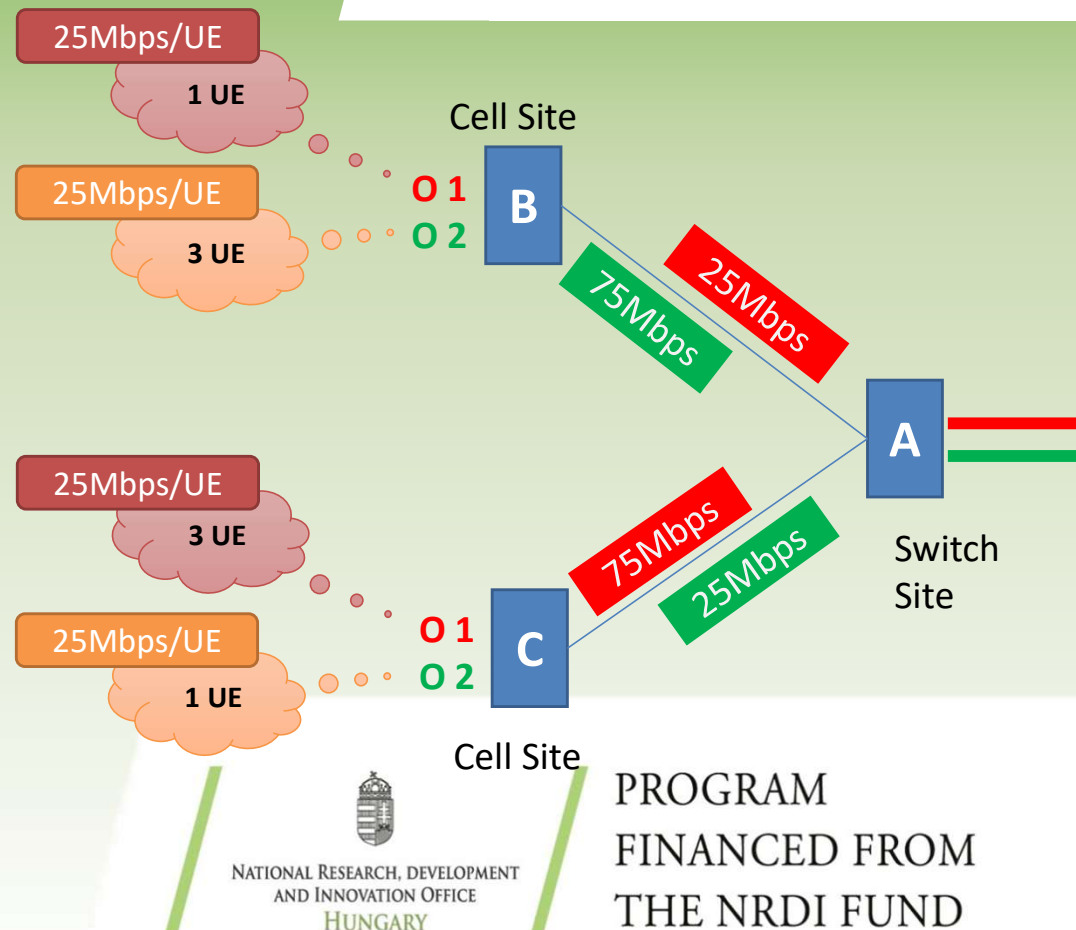
Case 2 „balanced asymmetry” Link-by-link HQoS

- Each link is shared according to the policy defined between the operators (1:1 in this example)
 - 100 Mbps allocated to both ops.
- The cell loads are not considered
 - The operator throughput shared evenly among them
- Each cell splits the given 50Mbps among the Ues within



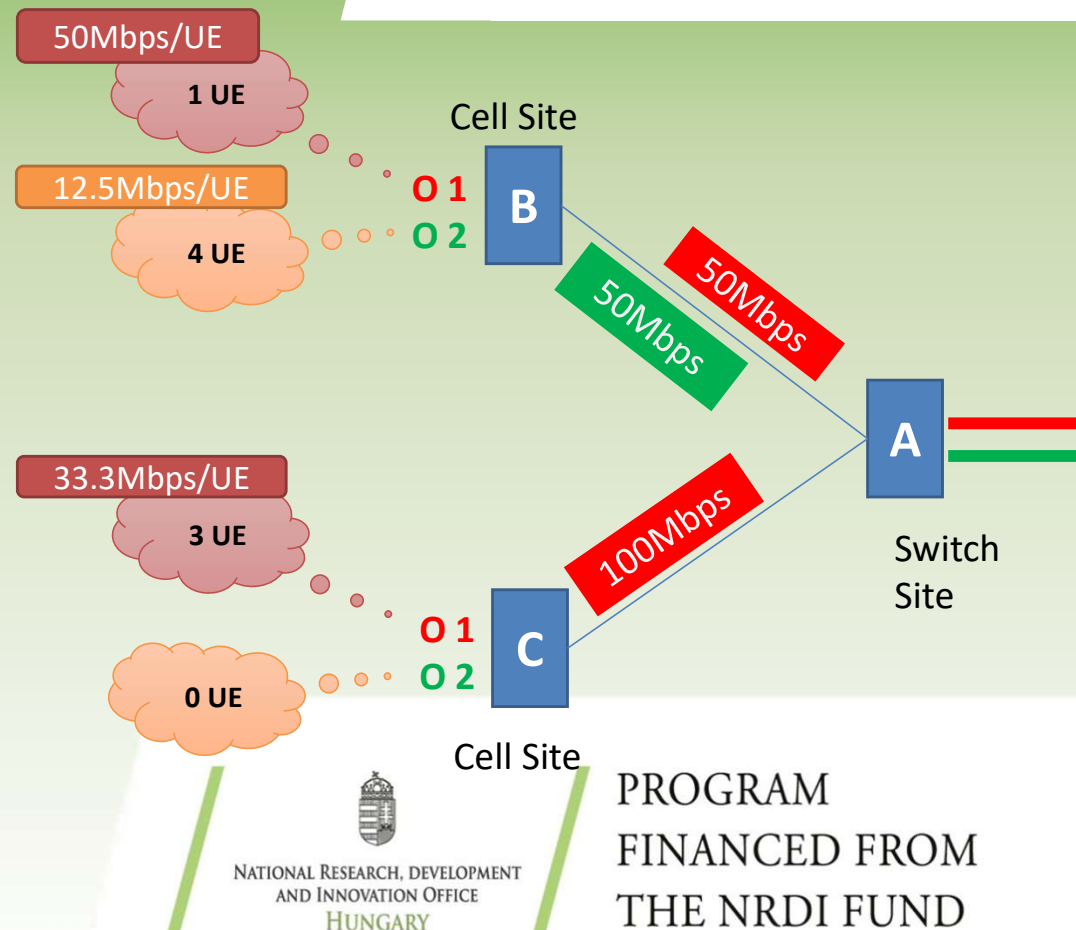
Case 2 „balanced asymmetry” Network-wide HQoS

- Each link is shared proportionally
 - number of users of each operator
 - operator shares according to the policy
- Considers cell load
- Better fairness among UEs



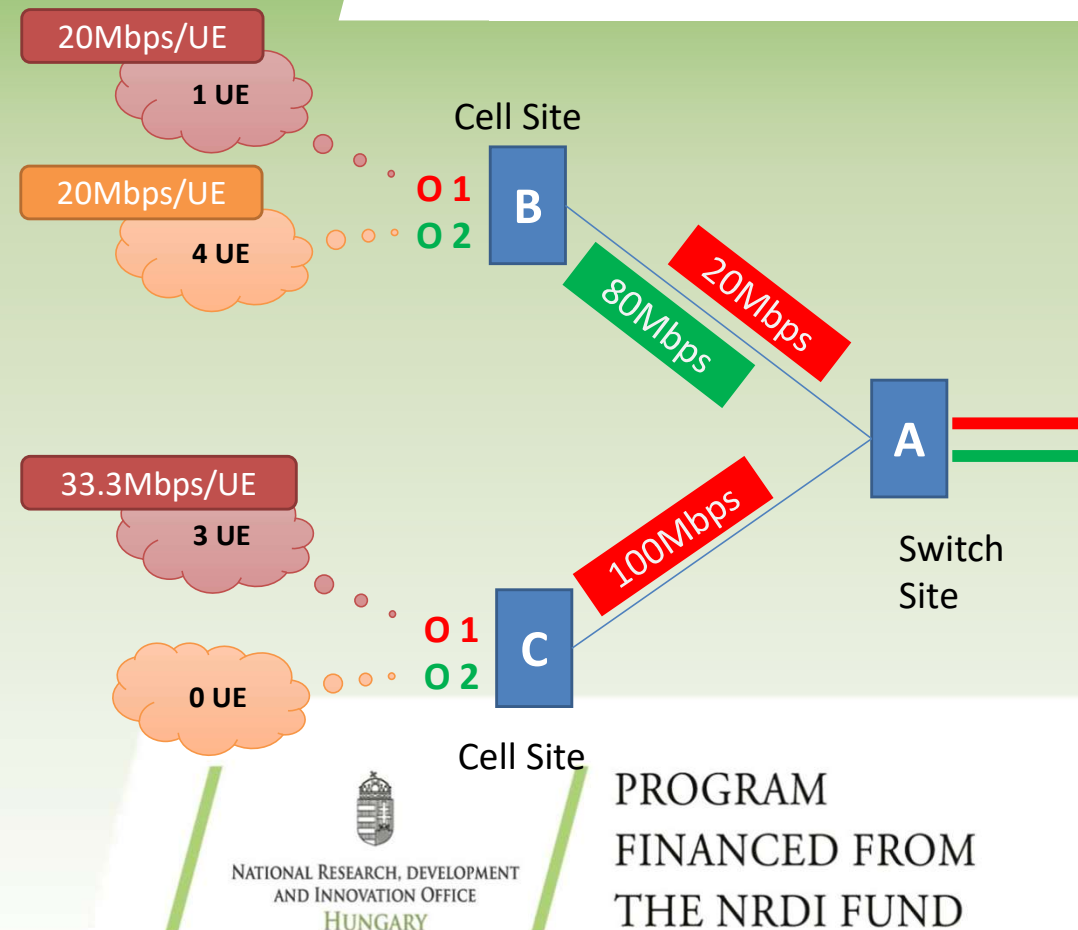
Case 4 „imbalanced asymmetry” Link-by-link HQoS

- Each link is shared according to the policy defined between the operators (1:1 in this example)
 - 100 Mbps allocated to both ops.
- Throughput allocated to UEs varies in a wide range 12.5 to 50Mbps
- Desired share between O1 and O2 is not met
 - 3:1 instead of 1:1



Case 4 „imbalanced asymmetry” Network-wide QoS

- The most congested link is first shared proportionally
 - to the operator weights 1:1 and the number of TAs (UEs) using the link
- On the most congested link the throughput is shared fairly among UEs in cell B
 - The other link is only used by O1 so UEs in cell C can get more without affecting UEs in cell B
- Desired share between O1 and O2 is much better than in link-by-link case
 - 12:8 instead of 1:1
 - 1:1 could only be achieved by starving a UE in cell B





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Thank you!

<http://ppv.elte.hu>



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